## On Computing Straight Skeletons by Means of Kinetic Triangulations

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#### Problem

Given a planar, straight line graph, construct the *straight skeleton*.

- Introduced by Aichholzer et al. in 1995.
- A skeleton consisting exclusively of straight line segments.
- Defined by a *wavefront propagation process*: The straight skeleton is the set of loci that are traced out by wavefront vertices.



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## Wavefront propagation

During this wavefront propagation, the wavefront topology changes:

- Split events: a wavefront vertex crashes into an oncoming wavefront edge.
- Edge events: a wavefront edge vanishes.



Extend from simple polygons to planar straight line graphs:



Extend from simple polygons/to planar straight line graphs:

# Applications







Tool path generation

Roof construction

Cut-and-fold

and more ...

- Keep a triangulation of the area not yet swept over by the wavefront [Aichholzer, Aurenhammer 1998].
- Edge and Split events are witnessed by collapsing triangles ⇒ Priority Queue.



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- These cases need special handling nevertheless. They are *flip events*.



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- We have implemented this algorithm, filling in a few gaps in the algorithm, including issues that arise from not assuming general position.
- We have run extensive tests using this code, more on that in a bit.

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- Unfortuantely the convex hull changes with time, and it matters.



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- But what if these edges are parallel?
- $\Rightarrow$  infinitely fast moving vertex.



















• Without general position, this algorithm can end up in infinite loops.



• This is not a result of inexact floating point operatons but can also happen with exact arithmetic.

- Keep a history of flip events  $\langle e_1, e_2, \ldots \rangle$ , where each  $e_i = (t_i, \Delta_i)$ .
- This history can be cleared when we encounter an edge or split event.
- If we encounter a flip event a second time, we may be in a flip event loop.

# Handling flip event loops

Brief outline:

- *e<sub>a</sub>* [*e<sub>b</sub>*] is the first [last] occurance of the duplicate event.
- Events between *e<sub>a</sub>* and *e<sub>b</sub>* happen at the same time.
- The set of triangles with events between *e<sub>a</sub>* and *e<sub>b</sub>* make up one or more edge-connected components.
- The component that contains the triangle of *e<sub>a</sub>* is a polygon *P* which has collapsed to a straight line.
- Undo the events of the triangles in *P*, and retriangulate *P* and its neighborhood.



 This approach also is applicable to kinetic triangulations in other algorithms.

#### Number of flip events

- Three points moving at constant speed become collinear at most twice.
- With *n* vertices, there are <sup>(n</sup><sub>3</sub>) ∈ O(n<sup>3</sup>) combinatorially different triangles.
- \$\mathcal{O}(n^3)\$ is the best known upper bound on the number of flip events!
- No input is known that results in more than quadratically many flip events.
- It turns out that for *practical data* the number of flip events is very linear.

#### Performance observations

	theoretical worst case		practical	
	runtime	space	runtime	space
E&E <sup>1</sup>	$\mathcal{O}(n^{17/11+\epsilon})$	$\mathcal{O}(n^{17/11+\epsilon})$	N/A	
CGAL <sup>2</sup>	$\mathcal{O}(n^2 \log n)$	$\mathcal{O}(n^2)$	$\mathcal{O}(n^2 \log n)$	$\mathcal{O}(n^2)$
Bone <sup>3</sup>	$\mathcal{O}(n^2 \log n)$	$\mathcal{O}(n)$	$\mathcal{O}(n \log n)$	$\mathcal{O}(n)$
Surfer <sup>4</sup>	$\mathcal{O}(n^3 \log n)$	$\mathcal{O}(n)$	$\mathcal{O}(n \log n)$	$\mathcal{O}(n)$

- <sup>1</sup>Eppstein and Erickson, 1999
- <sup>2</sup>F. Cacciola, 2004
- <sup>3</sup>Huber and Held, 2010
- <sup>4</sup>this, based on Aichholzer and Aurenhammer, 1998

#### Runtime tests



Runtime and memory usage behavior of CGAL, Bone, and Surfer for inputs of different sizes.

Bone and Surfer use their IEEE 754 double precision backend.

- We have implemented Aichholzer and Aurenhammer's algorithm from 1998, filling in details in the algorithm description. We fixed real problems that arise in the absence of general position.
- Our approach to handling flip events has wider applications.
- The implementation runs in  $\mathcal{O}(n \log n)$  time for *real-world data*. The number of flip events is linear in practice.
- It is industrial-strength, having been tested on tens of thousands of inputs.
- It is the fastest straight skeleton construction code to date, handling millions of vertices in mere seconds.

